

# On the VC-dimension of Uniform Hypergraphs

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## Abstract

Let  $\mathcal{F}$  be a  $k$ -uniform hypergraph on  $[n]$  where  $k - 1$  is a power of some prime  $p$  and  $n \geq n_0(k)$ . Our main result says that if  $|\mathcal{F}| > \binom{n}{k-1} - \log_p n + k!k^k$ , then there exists  $E_0 \in \mathcal{F}$  such that  $\{E \cap E_0 : E \in \mathcal{F}\}$  contains all subsets of  $E_0$ . This improves a longstanding bound of  $\binom{n}{k-1}$  due to Frankl and Pach [7].

## 1. Introduction

Let  $G$  be a set system (or hypergraph) on  $X$  and  $S$  be a subset of  $X$ . The *trace* of  $G$  on  $S$  is defined as  $G|_S = \{E \cap S : E \in G\}$ . We treat  $G|_S$  as a set and therefore omit multiplicity. We say that  $S$  is *shattered* by  $G$  if  $G|_S = 2^S$ , the set of all subsets of  $S$ . The Vapnik-Chervonenkis dimension (*VC dimension*) of  $G$  is the maximum size of a set shattered by  $G$ . Extremal problems on traces started from determining the maximum size of a set system on  $n$  vertices with VC dimension  $k - 1$  (equivalently, without a shattered  $k$ -set). Sauer [10], Perles and Shelah [11], and Vapnik and Chervonenkis [12] independently proved that this maximum is  $\binom{n}{0} + \dots + \binom{n}{k-1}$ . This and other results on traces have found numerous applications in geometry and computational learning theory (see Füredi and Pach [9] and Section 7.4 Babai and Frankl [3]).

Given two set systems  $G$  and  $F$ , if there exists a set  $S$  such that  $G|_S$  contains a copy of  $F$  as a subhypergraph, we say that  $G$  contains  $F$  as a trace. In this case we write  $G \rightarrow F$  ( $G \not\rightarrow F$  otherwise). Let  $\binom{X}{r}$  denote the set of all  $r$ -subsets of  $X$ . We call  $G$

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\*Research supported in part by NSF grants DMS-0400812 and an Alfred P. Sloan Research Fellowship.

†Research supported in part by NSA grant H98230-05-1-0079. Part of this research was done while working at University of Illinois at Chicago.

an  $r$ -uniform hypergraph ( $r$ -graph) on  $X$  if  $G \subseteq \binom{X}{r}$  and call the members of  $G$  edges. We define  $\text{Tr}^r(n, F)$  as the maximum number of edges in an  $r$ -graph on  $[n] = \{1, \dots, n\}$  not containing  $F$  as a trace. Frankl and Pach [7] considered the maximum size of uniform hypergraphs with fixed VC dimension. They showed that  $\text{Tr}^r(n, 2^{[k]}) \leq \binom{n}{k-1}$  for  $k \leq r \leq n$ . They conjectured that  $\text{Tr}^k(n, 2^{[k]}) = \binom{n-1}{k-1}$  for sufficiently large  $n$ . Obviously if a  $k$ -graph  $G$  contains a shattered edge, then  $G$  contains two disjoint edges (since the empty set appears in the trace). Therefore the conjecture of Frankl and Pach, if true, generalizes the well-known Erdős-Ko-Rado Theorem [5]. However, Ahlswede and Khachatrian [1] disproved it by constructing a  $G \subseteq \binom{[n]}{k}$  of size  $\binom{n-1}{k-1} + \binom{n-4}{k-3}$  that contains no shattered  $k$ -set when  $k \geq 3$  and  $n \geq 2k$ . Combining this with the upper bound in [7], for  $k \geq 3$  and  $n \geq 2k$ ,

$$\binom{n-1}{k-1} + \binom{n-4}{k-3} \leq \text{Tr}^k(n, 2^{[k]}) \leq \binom{n}{k-1}. \quad (1)$$

Our main result improves the upper bound in (1) in the case that  $k-1$  is a prime power and  $n$  is large.

**Theorem 1.** *Let  $p$  be a prime,  $t$  be a positive integer,  $k = p^t + 1$ , and  $n \geq n_0(k)$ . If  $\mathcal{F}$  is a  $k$ -uniform hypergraph on  $[n]$  with more than  $\binom{n}{k-1} - \log_p n + k!k^k$  edges, then there is a  $k$ -set shattered by  $\mathcal{F}$ . In other words,*

$$\text{Tr}^k(n, 2^{[k]}) \leq \binom{n}{k-1} - \log_p n + k!k^k.$$

In addition, we find exponentially many  $k$ -graphs achieving the lower bound in (1).

**Proposition 2.** *Let  $P(n, r)$  denote the number of non-isomorphic  $r$ -graphs on  $[n]$ . Then for  $k \geq 3$ , there are at least  $P(n-4, k-1)/2$  non-isomorphic  $k$ -graphs  $\mathcal{F}$  on  $[n]$  such that  $|\mathcal{F}| = \binom{n-1}{k-1} + \binom{n-4}{k-3}$  and  $\mathcal{F} \not\supseteq 2^{[k]}$ .*

Note that the gap between the upper and lower bounds in (1) is  $\binom{n-1}{k-2} - \binom{n-4}{k-3}$ . Theorem 1 reduces this gap by essentially  $\log n$  for certain values of  $k$ . Though this improvement is small, the value of Theorem 1 is perhaps mainly in its proof – a mixture of algebraic and combinatorial arguments. The main tool in proving  $\text{Tr}^k(n, 2^{[k]}) \leq \binom{n}{k-1}$  in [7] is the so-called *higher-order inclusion matrix*, whose rows are labeled by edges of a hypergraph  $\mathcal{F} \subseteq \binom{[n]}{k}$ . It was shown that if  $\mathcal{F}$  contains no shattered  $k$ -sets, then the rows of this matrix are linearly independent. Consequently  $|\mathcal{F}|$ , the number of the rows, equals to the rank of the matrix, which is at most  $\binom{n}{k-1}$ . The main idea in proving Theorem 1 is to enlarge the inclusion matrix of  $\mathcal{F}$  by adding more rows such that the rows in the enlarged matrix are still linearly independent. The method of adding independent vectors (or functions) to a space has been used before, *e.g.*, on the two-distance problem by Blokhuis [4] and a proof of the Ray-Chaudhuri–Wilson Theorem by Alon, Babai and Suzuki [2].

In order to prove Theorem 1, we also need more combinatorial tools. In particular, the sunflower lemma of Erdős and Rado [6], which is used to prove Lemma 3 below. Note that Lemma 3 and Theorem 4 together prove Theorem 1. Let  $2^{[k]-} = 2^{[k]} \setminus \emptyset$ .

**Lemma 3.** For any  $k \leq n$ ,

$$\mathrm{Tr}^k(n, 2^{[k]}) \leq \mathrm{Tr}^k(n, 2^{[k]^-}) + k!k^k.$$

**Theorem 4.** Let  $p$  be a prime,  $t$  be a positive integer, and  $k = p^t + 1$ . Then  $\mathrm{Tr}^k(n, 2^{[k]^-}) \leq \binom{n}{k-1} - \log_p n$  for  $n \geq n_0(k)$ .

In next section we prove Proposition 2 and Lemma 3. We prove Theorem 4 in Section 3 and give concluding remarks in the last section.

## 2. Proofs of Proposition 2 and Lemma 3.

**Proof of Proposition 2.** We construct  $\mathcal{F} = \mathcal{F}_0 \cup \mathcal{F}_1 \cup \mathcal{F}_2$  such that  $\mathcal{F}_0$  is the set of all  $k$ -sets containing 1 and 2, edges in  $\mathcal{F}_1$  contain 1 but avoid 2, and edges in  $\mathcal{F}_2$  contain 2 but avoid 1. If we let  $G_i = \{E \setminus \{i\} : E \in \mathcal{F}_i\}$  denote the link graph of  $i$  in  $\mathcal{F}_i$ , then  $G_1$  and  $G_2$  are  $(k-1)$ -graphs on  $V' = \{3, 4, \dots, n\}$ . Let  $G_1$  and  $G_2$  further satisfy the following conditions:

1.  $G_1 \cup G_2 = \binom{V'}{k-1}$
2.  $G_1 \cap G_2 = \{E \in \binom{V'}{k-1} : E \supseteq \{3, 4\}\}$
3.  $G_1 \supseteq \{E \in \binom{V'}{k-1} : E \ni 3, E \not\ni 4\}$ ,  $G_2 \supseteq \{E \in \binom{V'}{k-1} : E \ni 4, E \not\ni 3\}$ .

It is easy to see that  $|\mathcal{F}| = \binom{n-1}{k-1} + \binom{n-4}{k-3}$ , since  $|\mathcal{F}_0| = \binom{n-2}{k-2}$  and

$$|\mathcal{F}_1| + |\mathcal{F}_2| = |G_1| + |G_2| = |G_1 \cup G_2| + |G_1 \cap G_2| = \binom{n-2}{k-1} + \binom{n-4}{k-3}.$$

We claim that  $\mathcal{F} \not\cong 2^{[k]}$ . Suppose to the contrary that some  $E \in \binom{[n]}{k}$  is shattered. Then  $E \in \mathcal{F}$ . Note that every edge in  $\mathcal{F}$  contains either 1 or 2. If  $\{1, 2\} \subset E$ , then  $E \setminus \{1, 2\}$  is not contained in  $\mathcal{F}|_E$ . Without loss of generality, assume that  $E \ni 1$  and  $E \not\ni 2$ . Since  $E \setminus \{1\} \in G_1$  is contained in  $\mathcal{F}|_E$ , we have  $(E \setminus \{1\}) \cup \{2\} \in \mathcal{F}$  and consequently  $E \setminus \{1\} \in G_1 \cap G_2$ . Therefore  $E \supseteq \{3, 4\}$ . In order to have  $E \setminus \{1, 4\} \in \mathcal{F}|_E$ , there must be one edge of  $G_2$  containing 3 and not containing 4. But this is impossible because of the third condition on  $G_1$  and  $G_2$ .

In the above construction, every  $E \in \binom{V'}{k-1}$  with  $E \not\ni 3, E \not\ni 4$  could be in either  $G_1$  or  $G_2$ . These *undecided* edges form a complete  $(k-1)$ -graph  $K_{n-4}^{k-1}$  on  $\{5, \dots, n\}$ . Recall that  $P(n-4, k-1)$  is the number of non-isomorphic  $(k-1)$ -graphs on  $n-4$  vertices, or the number of non-isomorphic 2-edge-colorings of  $K_{n-4}^{k-1}$ . We claim that the number of non-isomorphic  $\mathcal{F}$  satisfying our construction is  $P(n-4, k-1)/2$ . To see this, let us consider

vertex degrees in  $\mathcal{F}$ . Let  $\deg(x)$  be the number of edges in  $\mathcal{F}$  containing a vertex  $x$ . It is not hard to see that no matter what the undecided edges are,  $\deg(1)$  and  $\deg(2)$  are always greater than  $\deg(3) = \deg(4)$ , which is greater than  $\deg(x)$  for all  $x > 4$ , and  $\deg(x)$  is fixed for all  $x > 4$ . Therefore two constructions  $\mathcal{F}$  and  $\mathcal{F}'$  are isomorphic if and only if  $\mathcal{F}|_{\{5, \dots, n\}}$  and  $\mathcal{F}'|_{\{5, \dots, n\}}$  are isomorphic or one is the complement of the other (since the vertices 1 and 2 are identical).  $\square$

Note that the construction in [1] is isomorphic to the case when all undecided  $E$  are in  $G_1$ .

A *sunflower* (or  $\Delta$ -*system*) with  $r$  petals and a core  $C$  is a collection of distinct sets  $S_1, \dots, S_r$  such that  $S_i \cap S_j = C$  for all  $i \neq j$ . Erdős and Rado [6] proved the following simple but extremely useful and fundamental lemma.

**Lemma 5 (Sunflower Lemma).** *Let  $G$  be a  $k$ -graph with  $|G| > k!(r-1)^k$ . Then  $G$  contains a sunflower with  $r$  petals.*  $\square$

We call a set  $S$  almost-shattered by  $\mathcal{F}$  if  $\mathcal{F}|_S$  contains  $2^S \setminus \emptyset$ .

**Proof of Lemma 3.** Let  $\mathcal{F}$  be a  $k$ -graph on  $[n]$  with  $|\mathcal{F}| > \text{Tr}^k(n, 2^{[k]-}) + k!k^k$ . We need to show that  $\mathcal{F}$  contains a shattered set. Since  $|\mathcal{F}| > \text{Tr}^k(n, 2^{[k]-})$ , we may find an almost-shattered  $k$ -set  $E_1 \in \mathcal{F}$ . Since  $|\mathcal{F} \setminus \{E_1\}| > \text{Tr}^k(n, 2^{[k]-})$ , we may find an almost-shattered  $k$ -set  $E_2 \in \mathcal{F} \setminus \{E_1\}$ . Repeating this process, we find distinct almost-shattered sets  $E_1, E_2, \dots, E_{k!k^k} \in \mathcal{F}$ . By the Sunflower Lemma,  $\mathcal{F}' = \{E_1, \dots, E_{k!k^k}\}$  contains a sunflower with  $k+1$  petals. Let us simply denote it by  $E_1, \dots, E_{k+1}$  and  $C = \bigcap_{i=1}^{k+1} E_i$ . Since  $E_1$  is almost-shattered by  $\mathcal{F}$  and  $E_1 \setminus C \neq \emptyset$ , there is  $E_0 \in \mathcal{F}$  such that  $E_0 \cap E_1 = E_1 \setminus C$ . Now  $E_1 \cap E_0, E_2 \cap E_0, \dots, E_{k+1} \cap E_0$  are pairwise disjoint. Since  $|E_0| = k < k+1$ , there exists  $i \neq 1$  such that  $E_i \cap E_0 = \emptyset$ . This means that  $\emptyset \in \mathcal{F}|_{E_i}$ . Consequently  $E_i$  is shattered by  $\mathcal{F}$ .  $\square$

### 3. Proof of Theorem 4

#### 3.1. Inclusion Matrices and Proof Outline

The proof of Theorem 4 needs the concept of higher-order inclusion matrices. Let  $\mathcal{F}$  be a set system on  $X$ . The *incidence matrix*  $M(\mathcal{F}, \leq s)$  of  $\mathcal{F}$  over  $\binom{X}{\leq s}$  is the matrix whose rows (*incidence vectors*) are labeled by the edges of  $\mathcal{F}$ , columns are labeled by subsets of  $[n]$  of size at most  $s$ , and entry  $(E, S)$ ,  $E \in \mathcal{F}$ ,  $|S| \leq s$ , is 1 if  $S \subseteq E$  and 0 otherwise. Throughout this paper, we fix  $s = k-1$  and simply write  $M(\mathcal{F})$  instead of  $M(\mathcal{F}, \leq k-1)$ . In particular, let

$$I(k) = M\left(\binom{[n]}{k}\right) = M\left(\binom{[n]}{k}, \leq k-1\right).$$

For each  $E \subset [n]$ , the incidence vector  $v_E$  is a  $(0, 1)$ -vector of length  $\binom{n}{0} + \dots + \binom{n}{k-1}$ , whose coordinates are labeled by all subsets of  $[n]$  of size at most  $k-1$ . Note that  $v_E$  always has a 1 in the position corresponding to  $\emptyset$ . Let  $e_i = v_{\{i\}}$  for each  $i \in [n]$ .

Let  $q$  be 0 or a prime number. As usual,  $\mathbb{F}_q$  denotes a field of  $q$  elements when  $q$  is a prime. Let us define  $\mathbb{F}_0$  to be  $\mathbb{Q}$ , the field of rational numbers. Given a hypergraph  $\mathcal{F}$ , a *weight function* of  $\mathcal{F}$  over  $\mathbb{F}_q$  is a function  $\alpha : \mathcal{F} \rightarrow \mathbb{F}_q$ . If  $\alpha(E) = 0$  for all  $E \in \mathcal{F}$ , then we call  $\alpha$  the zero function and write  $\alpha \equiv 0$ . We define

$$v(\mathcal{F}, \alpha) = \sum_{E \in \mathcal{F}} \alpha(E) v_E$$

and write  $v(\mathcal{F}) = \sum_{E \in \mathcal{F}} v_E$ . We say that  $\mathcal{F}$  is *linearly independent* in characteristic  $q$  if the rows of  $M(\mathcal{F})$  are linearly independent over  $\mathbb{F}_q$ , namely,  $v(\mathcal{F}, \alpha) = 0 \pmod{q}$  implies that  $\alpha \equiv 0$ .

Part 1 of Lemma 6 below is the key observation to the proof of the upper bound in (1). It implies that if  $\mathcal{F} \subseteq \binom{[n]}{k}$  contains no shattered sets, then it is linearly independent in any characteristic. Our proof of Theorem 4 also needs Part 2. We call a set  $S$  *near-shattered* by  $\mathcal{F}$  if  $\mathcal{F}|_S$  contains  $2^S \setminus (\{i\} \cup \emptyset)$  for some  $i \in S$ .

**Lemma 6.** *Let  $q$  be 0 or a prime number. Suppose that  $\mathcal{F} \subseteq \binom{[n]}{k}$  and  $\alpha : \mathcal{F} \rightarrow \mathbb{F}_q$  is a non-zero weight function. Define  $d(S) = \sum_{S \subseteq E \in \mathcal{F}} \alpha(E)$  for every subset  $S \subset [n]$ . Fix  $A \in \mathcal{F}$  with  $\alpha(A) \neq 0$ .*

1. *If  $d(S) = 0 \pmod{q}$  for all  $S \subset A$ , then  $A$  is shattered by  $\mathcal{F}$ .*
2. *Let  $i \in A$ . If  $d(S) = 0 \pmod{q}$  for all  $S \subset A$  with  $S \neq \emptyset$  and  $S \neq \{i\}$ , then  $A$  is near-shattered.*

**Proof.** Part 1 and Part 2 have almost the same proofs. Since Part 1 was proved in [7] and [3] (Theorem 7.27), we only prove Part 2 here.

Since  $\mathcal{F}$  is  $k$ -uniform, we have  $d(A) = \alpha(A) \neq 0$ . For  $B \subseteq A$ , we define  $d(A, B) = \sum_{E \in \mathcal{F}, E \cap A = B} \alpha(E)$ . The following equality can be considered as a variant of the Inclusion-Exclusion formula.

$$d(A, B) = \sum_{B \subseteq S \subseteq A} (-1)^{|S-B|} d(S). \quad (2)$$

In fact, because  $d(B) = d(A, B) + \sum_{E \in \mathcal{F}, B \subseteq E \cap A} \alpha(E)$ , (2) is equivalent to

$$\sum_{E \in \mathcal{F}, B \subseteq E \cap A} \alpha(E) + \sum_{B \subseteq S \subseteq A} (-1)^{|S-B|} d(S) = 0.$$

This holds because on the left side, each  $\alpha(E)$  with  $r = |E \cap A| - |B| > 0$  has coefficient  $1 - \binom{r}{1} + \dots + (-1)^r \binom{r}{r} = 0$ .

Pick any  $B \subset A$  with  $B \neq \emptyset$  and  $B \neq \{i\}$ . We now show that there exists  $E \in \mathcal{F}$  such that  $E \cap A = B$ . We use (2) and the assumption that  $d(S) = 0 \pmod{q}$  for all  $S$  with  $B \subseteq S \subset A$  to derive

$$\sum_{E \in \mathcal{F}, E \cap A = B} \alpha(E) = d(A, B) = \sum_{B \subseteq S \subseteq A} (-1)^{|S-B|} d(S) = (-1)^{|A-B|} d(A) \neq 0 \pmod{q}.$$

Hence the sum on the left side is not empty.  $\square$

By Lemma 6 Part 1, if  $\mathcal{F}$  contains no shattered sets, then the rows of  $M(\mathcal{F})$  are linearly independent (over  $\mathbb{Q}$ ) and consequently  $|\mathcal{F}| = \text{rank}(M)$ . Clearly  $\text{rank}(M) \leq \text{rank}(I(k))$ . It is well-known that  $\text{rank}_{\mathbb{Q}}(I(k)) = \binom{n}{k-1}$  (e.g., see [3] section 7.3). This immediately gives  $\text{Tr}^k(n, 2^{[k]}) \leq \binom{n}{k-1}$ , the result of Frankl and Pach [7].

The proof of Theorem 4 proceeds as follows. Suppose that  $\mathcal{F} \subseteq \binom{[n]}{k}$  satisfies  $\mathcal{F} \not\prec 2^{[k]-}$ . Recall that  $k = p^t + 1$  for some prime  $p$  and positive integer  $t$ . We will construct a matrix  $M'$  obtained from  $M = M(\mathcal{F})$  by adding  $\log_p n$  new rows. The new rows have the form  $e_S = \sum_{i \in S} e_i$ , for some set  $S$  of size  $m = p^{t+1}$ . In other words, a new row has entry 1 at  $m$  coordinates corresponding to  $m$  singletons and 0 otherwise (the entry at  $\emptyset$  is 0 because  $m = 0 \pmod p$ ). The main step is to show that these new rows lie in the row space of  $I(k)$ , and all the rows of  $M'$  are still linearly independent. Consequently,

$$|\mathcal{F}| + \log_p n = \text{rank}_{\mathbb{F}_p}(M') \leq \text{rank}_{\mathbb{F}_p}(I(k)) \leq \text{rank}_{\mathbb{Q}}(I(k)) = \binom{n}{k-1},$$

which implies that  $|\mathcal{F}| \leq \binom{n}{k-1} - \log_p n$ .

We now divide the main step into three lemmas, which we will prove in the next subsection.

**Lemma 7.** *Suppose that  $k = p^t + 1$  and  $m = p^{t+1}$  for prime  $p$  and  $t > 0$ . Then for every  $S \in \binom{[n]}{m}$ ,  $e_S$  is in the row space of  $I(k)$  over  $\mathbb{F}_p$ .*

Lemma 8 is the key to our proof. For  $a, b \in [n]$ , let  $e_{a,-b} = e_a - e_b$ . Thus  $e_{a,-b}$  is the vector with a 1 in position  $\{a\}$ , a  $-1$  in position  $\{b\}$ , and 0 everywhere else. Lemma 8 says that  $e_{a,-b}$  is outside the row space of  $M$  for every  $a \neq b$ .

**Lemma 8.** *Let  $k \geq 2$  and  $n \geq n_0(k)$ . Suppose that  $\mathcal{F} \subseteq \binom{[n]}{k}$  contains no almost-shattered set, i.e.,  $\mathcal{F} \not\prec 2^{[k]-}$ . If  $|\mathcal{F}| > \binom{n}{k-1} - \log_p n$ , then for every two distinct  $a, b \in [n]$ , the set  $\{v_E : E \in \mathcal{F}\} \cup \{e_{a,-b}\}$  is linearly independent in any characteristic.*

**Lemma 9.** *Given a prime  $p$  and  $m \geq 1$ , let  $n \geq n_0(p, m)$  and  $r = \log_p n$ . Suppose that for every two distinct  $a, b \in [n]$ , the set  $\{v_E : E \in \mathcal{F}\} \cup \{e_{a,-b}\}$  is linearly independent in characteristic  $p$ . Then there exist subsets  $S_1, \dots, S_r \in \binom{[n]}{m}$  such that the set  $\{v_E : E \in \mathcal{F}\} \cup \{e_{S_1}, \dots, e_{S_r}\}$  is linearly independent in characteristic  $p$ .*

### 3.2. Proof of Lemmas

Given a hypergraph  $\mathcal{F}$  on  $X$  and a subset  $A \subseteq X$ , we define the *degree*  $\deg_{\mathcal{F}}(A)$  to be the number of edges in  $\mathcal{F}$  containing  $A$ .

**Proof of Lemma 7.** Let  $K = \binom{[n]}{k}$ . It suffices to prove that  $\sum_{E \in K} v_E = c \cdot e_S$  for some nonzero  $c \in \mathbb{F}_p$ . Equivalently, we need to show that for  $T \subset S$ ,  $\deg_K(T) = 0 \pmod p$  when

$|T| \geq 2$  or  $|T| = 0$ , and  $\deg(T) = c \neq 0 \pmod p$  when  $|T| = 1$ . Since  $K$  is a complete  $k$ -graph,  $\deg_K(T) = \binom{m-|T|}{k-|T|}$ . By a well-known result of Kummer, the binomial coefficient  $\binom{a}{b}$  is divisible by a prime  $p$  if and only if, when writing  $a$  and  $b$  as two numbers in base  $p$ ,  $a = (a_j \cdots a_1 a_0)_p$  and  $b = (b_j \cdots b_1 b_0)_p$ , there exists  $i \leq j$ , such that  $b_i > a_i$ . Since  $m$  is a power of  $p$ , for any  $1 \leq k \leq m-1$ ,  $p$  divides  $\binom{m}{k}$ . Hence  $\deg_K(\emptyset) = \binom{m}{k} = 0 \pmod p$ . Now consider  $|T| = s \geq 2$ . Since  $k = p^t + 1$ , we know  $k-s < p^t$  and thus write  $k-s = (a_{t-1} \dots a_0)_p$ . Since  $m = p^{t+1}$ , we have  $m-s = p^{t+1} - s = (p-1)p^t + k-s-1$ . We thus have  $m-s = (p-1 a_{t-1} \dots a_0)_p - 1$ . Hence there exists  $i \leq t-1$  such that the value of  $m-s$  at bit  $i$  is less than  $a_i$  and consequently  $\binom{m-s}{k-s}$  is divisible by  $p$ . When  $|T| = 1$ , we have  $m-1 = p^{t+1} - 1$  and therefore  $\binom{m-1}{k-1}$  is not divisible by  $p$  for any  $1 \leq k \leq m-1$ .  $\square$

**Proof of Lemma 8.** We prove the contrapositive of the claim: If  $\mathcal{F} \not\rightarrow 2^{[k]-}$  and there exists a non-zero function  $\alpha : \mathcal{F} \rightarrow \mathbb{F}_q$  such that  $v(\mathcal{F}, \alpha) = e_{a,-b}$  for some  $a, b \in [n]$  ( $a \neq b$ ), then  $|\mathcal{F}| \leq \binom{n}{k-1} - \log_p n$ . We claim that it suffices to show that  $\deg_{\mathcal{F}}(\{a\}) = O(n^{k-3})$ . In fact, suppose  $\deg_{\mathcal{F}}(\{a\}) \leq c_k n^{k-3}$  for some constant  $c_k$  and  $|\mathcal{F}| > \binom{n}{k-1} - \log_p n$ . After we remove  $a$  and all the edges containing  $a$ , we obtain a  $k$ -graph  $\tilde{\mathcal{F}} \subseteq \mathcal{F}$  with  $n-1$  vertices satisfying

$$\begin{aligned} |\tilde{\mathcal{F}}| &> \binom{n}{k-1} - \log_p n - c_k n^{k-3} \\ &= \binom{n-1}{k-1} + \binom{n-1}{k-2} - \log_p n - c_k n^{k-3} \\ &\geq \binom{n-1}{k-1} \end{aligned}$$

where the last inequality holds because  $\binom{n-1}{k-2} \geq \log_p n + c_k n^{k-3}$  for  $n \geq n_0(k)$ . But we showed that  $\text{Tr}^k(n, 2^{[k]}) \leq \binom{n}{k-1}$  for any  $k \leq n$ , therefore  $\tilde{\mathcal{F}} \rightarrow 2^{[k]}$ , a contradiction.

Suppose that  $\sum_{E \in \mathcal{F}} \alpha(E) v_E = e_{a,-b}$ . Let  $\mathcal{F}' = \{E \in \mathcal{F} : \alpha(E) \neq 0\}$  and  $V' = [n] \setminus \{a, b\}$ . For a subset  $A \subset [n]$ , let  $d(A) = \sum_{A \subseteq E \in \mathcal{F}'} \alpha(E) \pmod q$ . Our assumption  $v(\mathcal{F}, \alpha) = e_{a,-b}$  implies that  $d(\{a\}) = 1$ ,  $d(\{b\}) = -1$ , and  $d(A) = 0$  for every  $A \neq \{a\}, \{b\}$  and  $|A| \leq k-1$ . Applying Lemma 6 Part 1, we conclude that no  $E \in \mathcal{F}'$  satisfies  $E \subseteq V'$ . In other words, every edge in  $\mathcal{F}'$  contains either  $a$  or  $b$ . Next observe that if  $\mathcal{F}'$  contains an edge  $E$  such that  $a \in E$  and  $b \notin E$ , then  $\mathcal{F}'$  also contains  $(E \setminus \{a\}) \cup \{b\}$ . Otherwise  $E$  is the only edge in  $\mathcal{F}'$  containing  $E \setminus \{a\}$  and consequently  $d(E \setminus \{a\}) = \alpha(E) \neq 0$ , a contradiction.

Let  $G_a = \{E \setminus \{a\} : E \in \mathcal{F}', a \in E, b \notin E\}$  and define  $G_b$  similarly. By the previous observation, we have  $G := G_a = G_b$ . We then observe that  $G \neq \emptyset$  otherwise every edge (of  $\mathcal{F}'$ ) containing  $a$  also contains  $b$ , and consequently  $1 = d(\{a\}) = d(\{a, b\}) = 0$ .

Fix an edge  $E_0 \in \mathcal{F}'$  containing  $a$  but not  $b$ . Applying Lemma 6 Part 2, we conclude that  $E_0$  is near-shattered, *i.e.*, all subsets of  $E_0$  are in the trace  $\mathcal{F}'|_{E_0}$  except for  $\{a\}$  and  $\emptyset$ . If another edge  $E \in \mathcal{F}$  satisfies  $E \cap E_0 = \{a\}$ , then  $E_0$  becomes almost-shattered, contradicting the assumption that  $\mathcal{F} \not\rightarrow 2^{[k]-}$ . We may therefore assume that every  $E \in \mathcal{F}$  containing

$a$  also contains some other element of  $E_0$ . Below we show that there exists  $H \subseteq G$  with at most  $2k$  vertices and transversal number at least 2 (*i.e.*, no element lies in all sets of  $H$ ). Therefore every  $E \in \mathcal{F}$  containing  $a$  has at least two vertices in  $H$  and consequently  $\deg_{\mathcal{F}}(\{a\}) \leq \binom{2k}{2} \binom{n-3}{k-3} = O(n^{k-3})$ .

Pick  $A \in G_a$  (thus  $|A| = k - 1$ ). We claim that for every  $S \subset A$ ,  $|S| = k - 2$ , there exists  $B \in G_a$  such that  $A \cap B = S$ . Suppose instead, that for some  $S \in \binom{A}{k-2}$ , no such  $B$  exists. In this case,  $A \cup \{a\}$  and  $S \cup \{a, b\}$  are the only possible edges in  $\mathcal{F}'$  containing  $S \cup \{a\}$ . We thus have  $S \cup \{a, b\} \in \mathcal{F}'$ , otherwise  $d(S \cup \{a\}) = \alpha(A \cup \{a\}) \neq 0$ . Because  $G_a = G_b$ , no  $B \in G_b$  satisfies  $A \cap B = S$ . We now have a contradiction since

$$d(S) = \alpha(A \cup \{a\}) + \alpha(A \cup \{b\}) + \alpha(S \cup \{a, b\}) = d(A) + \alpha(S \cup \{a, b\}) = \alpha(S \cup \{a, b\}) \neq 0.$$

Now, for every  $S \in \binom{A}{k-2}$ , we choose exactly one set  $B = B(S) \in G_a$  such that  $A \cap B = S$ . Let  $H = \{A\} \cup \{B(S) : S \in \binom{A}{k-2}\}$ . Clearly  $H$  contains at most  $2k$  vertices. It is easy to see that there is no  $x \in \cap_{E \in H} E$ . In fact, if such  $x \in A$ , then  $B(A \setminus \{x\})$  misses  $x$ . If  $x \notin A$ , then  $A$  misses  $x$ . Therefore the transversal number of  $H$  is at least 2, and the proof is complete.  $\square$

**Proof of Lemma 9.** Let  $M$  be the inclusion matrix of  $\mathcal{F}$ . We sequentially add vectors  $e_{S_1}, \dots, e_{S_i}$  with  $S_1, \dots, S_i \in \binom{[n]}{m}$  to  $M$  such that  $e_{S_1}, \dots, e_{S_i}$  and the rows of  $M$  are linearly independent. We claim that this can be done as long as  $i \leq \log_p n$ . Suppose to the contrary, that there exists  $i \leq \log_p n - 1$  such that we fail to add a new vector at step  $i + 1$ . In other words, we have chosen  $e_{S_1}, \dots, e_{S_i}$  successfully, but for every  $S \in \binom{[n]}{m} \setminus \{S_1, \dots, S_i\}$ , there exist a weight function  $\alpha$  and  $c_1, \dots, c_i \in \mathbb{F}_p$  such that

$$e_S = v(\mathcal{F}, \alpha) + \sum_{j=1}^i c_j e_{S_j}. \quad (3)$$

We observe that for fixed  $c_1, \dots, c_i$ , the set of  $m$ -sets satisfying (3) forms a partial Steiner system  $PS(n, m, m - 1)$  (an  $m$ -graph on  $[n]$  such that each  $(m - 1)$ -subset of  $[n]$  is contained in at most one edge). In fact, if two  $m$ -sets  $S, S'$  with  $|S \cap S'| = m - 1$  both satisfy (3), with weight functions  $\alpha_1$  and  $\alpha_2$  respectively, then  $v(\mathcal{F}, \alpha_1 - \alpha_2) = e_{a, -b}$ , where  $\{a\} = S \setminus S'$  and  $\{b\} = S' \setminus S$ . This is a contradiction to our assumption. Consequently for fixed  $c_1, \dots, c_i$ , the number of  $m$ -sets satisfying (3) is at most  $\binom{n}{m-1}/m$ . As a result, the number of  $m$ -sets that cannot be chosen is at most  $p^i \binom{n}{m-1}/m$ . We thus obtain

$$\left| \binom{[n]}{m} \setminus \{S_1, \dots, S_i\} \right| = \binom{n}{m} - i \leq p^i \frac{1}{m} \binom{n}{m-1},$$

which implies that

$$(n - m + 1) - \frac{im}{\binom{n}{m-1}} \leq p^i.$$

Since  $i \leq \log_p n - 1$ , we have  $p^i \leq n/p$ , and consequently  $n - m + 1 - im/\binom{n}{m-1} \leq n/p$ , which is impossible for fixed  $p \geq 2, m$  and sufficiently large  $n$ .  $\square$

## 4. Concluding Remarks

We believe the lower bound in (1) is correct, though verifying this for all  $k$  may be hard because Proposition 2 gives exponentially many extremal hypergraphs. In order to reduce the bound in Theorem 1, one probably wants to look for a better way to find independent vectors than the greedy algorithm we used in the proof of Lemma 9. It may not be very hard to check this for the  $k = 3$  case, namely, to verify that  $\text{Tr}^3(n, 2^{[3]}) = \binom{n-1}{2} + 1$ . Using more involved combinatorial arguments, instead of the Sunflower Lemma, we can prove that  $\text{Tr}^3(n, 2^{[3]}) \leq \binom{n}{2} - \log_2 n$ .

Improving the upper bound further for other values of  $k$  will most likely need some new ideas. Our approach uses incidence vectors of a family of singletons. The following proposition shows that this approach requires  $k - 1$  to be a prime power.

**Proposition 10.** *Let  $p$  be a prime and  $k \geq 2$ . Suppose that  $\mathcal{F} \subseteq \binom{[n]}{k}$  and  $\alpha : \mathcal{F} \rightarrow \mathbb{F}_p$  is a non-zero weight function. Define  $d(S) = \sum_{S \subseteq E \in \mathcal{F}} \alpha(E)$  for every subset  $S \subset [n]$ . If there exists a vertex  $x \in [n]$  such that  $d(\{x\}) \neq 0$  and  $d(S) = 0 \pmod p$  for every  $S \ni x$  with  $2 \leq |S| \leq k - 1$ , then  $k - 1$  is a power of  $p$ .*

**Proof.** Let  $2 \leq s \leq k - 1$ . When we sum up  $d(S)$  for all  $S \ni x$  with  $|S| = s$ , we over-count  $d(\{x\})$  by a factor of  $\binom{k-1}{s-1}$ . In other words,

$$d(\{x\}) = \frac{1}{\binom{k-1}{s-1}} \sum_{x \in S, |S|=s} d(S).$$

Since  $d(\{x\}) \neq 0$  but  $d(S) = 0 \pmod p$  for all  $S$  in the right side, it must be the case that  $p$  divides  $\binom{k-1}{s-1}$ . We thus conclude that  $p$  divides  $\binom{k-1}{i}$  for all  $1 \leq i \leq k - 1$ . By the result of Kummer on binomial coefficients, this happens only if  $k - 1$  is a power of  $p$ .  $\square$

## 5. Acknowledgments

We thank the referees for their comments, which improved the presentation of the paper and in particular, shortened the proof of Lemma 9.

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